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Instructions: This examination paper includes 9 pages and 20 multiple-choice questions starting on page 3. You are responsible for ensuring that your copy of the paper is complete. Bring any discrepancy to the attention of your invigilator. The answers for all the questions must be indicated by filling the corresponding circle on the optical scanning (OMR) examination sheet. This OMR examination sheet is the only page to be handed in. The instructions for completing the OMR examination sheet are provided on page 2. Read and follow these instructions with care! There is one mark for each question. Answer all questions. There is no penalty for guessing. This is a closed book exam. No reference material of any kind is permitted. No calculators of any kind are permitted. Time allowed is $\mathbf{5 0}$ minutes.

Note: $\mathrm{A}^{\prime}$ and $\overline{\mathrm{A}}$ are used interchangeably.

## OMR examination instructions

## Multiple choice questions (numbered 1 to 20) - indicate your answer by filling the corresponding circle on the OMR answer sheet

1. The binary representation of $(-22)_{10}$ with 6 bits in 2 's complement format is:
2. $(010110)_{2}$
3. $\quad(110110)_{2}$
4. $(101010)_{2}$
5. $(101001)_{2}$
6. none of the above
7. The hexadecimal equivalent of the unsigned number represented as $(4444)_{8}$ in octal is:

| 1. | $(444)_{16}$ |
| :--- | :--- |
| 2. | $(814)_{16}$ |
| 3. | $(924)_{16}$ |
| 4. | $(\mathrm{A} 84)_{16}$ |
| 5. | $(\mathrm{B} 22)_{16}$ |

3. The largest positive number in the sign-magnitude format represented with 12 -bits is:
4. $(\mathrm{FFF})_{16}$
5. $(4096)_{10}$
6. $(1777)_{8}$
7. $(100000000000)_{2}$
8. none of the above
9. The range of numbers that can be represented with 9 bits in 1's complement format is:
10. $(-255)_{10}$ to $(255)_{10}$
11. $(-255)_{10}$ to $(256)_{10}$
12. $(-256)_{10}$ to $(255)_{10}$
13. $(-256)_{10}$ to $(256)_{10}$
14. $(0)_{10}$ to $(511)_{10}$
15. Consider the circuit from Figure 1. Function $F$ is:
16. $\quad F(A, B)=A B$
17. $F(A, B)=A^{\prime} B^{\prime}$
18. $\quad \mathbf{F}(\mathbf{A}, \mathbf{B})=A B^{\prime}$
19. $\quad \mathbf{F}(A, B)=A^{\prime} B$
20. none of the above


Figure 1 - Circuit for question 5.
6. Function $F(A, B, C, D)$ shown in Figure 2 is:

1. $\quad \Pi \mathbf{M}(0,5,10,15)$
2. $\quad \Pi \mathbf{M}(0,5,7,10,15)$
3. $\quad \Pi \mathbf{M}(1,2,3,4,6,8,9,11,12,13,14)$
4. $\quad \Pi \mathbf{M}(1,2,3,4,6,7,8,9,11,12,13,14)$
5. none of the above
6. All the prime implicants of $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ shown in Figure 2 are:
7. $A^{\prime} C, A C^{\prime}, A^{\prime} B^{\prime} C, A^{\prime} C D^{\prime}, B^{\prime} D, D^{\prime}$
8. $A C^{\prime}, A^{\prime} B^{\prime} C, A^{\prime} C D^{\prime}, B^{\prime} D, D^{\prime}$
9. $A C^{\prime}, A^{\prime} C D^{\prime}, B^{\prime} D, D^{\prime}$
10. $A C^{\prime}, A^{\prime} B^{\prime} C, B^{\prime} D, D^{\prime}$
11. A'C, A'B'C, B'D, BD'
12. All the essential prime implicants of $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ shown in Figure 2 are:
13. $A C^{\prime}, A^{\prime} B^{\prime} C, A^{\prime} C D^{\prime}, B^{\prime} D, B^{\prime}$
14. $A C^{\prime}, A^{\prime} B^{\prime} C, B^{\prime} D, B^{\prime}$
15. $A C^{\prime}, A^{\prime} C D^{\prime}, B^{\prime} D, D^{\prime}$
16. $A^{\prime}$ ', $B^{\prime} D, B^{\prime}$
17. $A^{\prime} C, B^{\prime} D, D^{\prime}$

| $C D$ | 00 | 01 | 11 | 10 |
| :---: | :---: | :---: | :---: | :---: |
| 00 | $0$ | 1 | 1 | 1 |
| 01 | 1 | 0 | 1 | 1 |
| 11 | 1 | 0 | 0 | 1 |
| 10 | 1 | 1 | 7 | 0 |

Figure 2 - Karnaugh map for function F(A,B,C,D) for questions 6 to 8.


Figure 3 - Karnaugh map for function F(A,B,C,D) for questions 9 and 10. Note, ' $X$ ' stands for don't care.
9. Consider the function $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ shown in Figure 3. The simplified logical expression in the sum-of-products (SOP) form (i.e., the minimum number of product terms and the minimum number of literals in every product term) for $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ can be converted into a circuit implementation using NAND gates, which is shown in:

1. Figure 4(a)
2. Figure 4(b)
3. Figure 4(c)
4. Figure 4(d)
5. all of the above

(a)

(b)

(c)

(d)

Figure 4 - Implementations for function $F(A, B, C, D)$ for question 9.
10. Consider the function $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ shown in Figure 3. The simplified logical expression in the product-of-sums (POS) form (i.e., the minimum number of sum terms and the minimum number of literals in every sum term) for $\mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C}, \mathbf{D})$ can be converted into a circuit implementation using NOR gates, which is shown in:

1. Figure 5(a)
2. Figure 5(b)
3. Figure 5(c)
4. Figure 5(d)
5. all of the above

(a)

(b)

(c)

(d)

Figure 5 - Implementations for function $F(A, B, C, D)$ for question 10.
11. The circuit from Figure 6 is equivalent to the circuit from:

1. Figure 7(a)
2. Figure 7(b)
3. Figure 7(c)
4. Figure 7(d)
5. none of the above



Figure 6 - Circuit for question 11.

(a)

(b)

(c)

(d)

Figure 7 - Circuits for question 11.
12. The circuit from Figure 8 is equivalent to the circuit from:

1. Figure 9(a)
2. Figure 9(b)
3. Figure 9(c)
4. Figure 9(d)
5. none of the above


Figure 8 - Circuit for question 12.

(a)

(b)

(c)

(d)

Figure 9-Circuits for question 12.


| $D_{2}$ | $D_{1}$ | $D_{0}$ | $Y_{0}$ | $Y_{1}$ | $Y_{2}$ | $Y_{3}$ | $Y_{4}$ | $Y_{5}$ | $Y_{6}$ | $Y_{7}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |

(a)


Figure 10 - The 3-to-8 decoder without an enable signal for questions 13 and 14.
13. In Figure 10(b):

1. $\quad \mathbf{F}(A, B, C)=\Pi M(3,5)$
2. $\quad F(A, B, C)=\Sigma \mathbf{m}(3,5,6)$
3. $\quad \mathbf{F}(\mathbf{A}, \mathrm{B}, \mathrm{C})=\Pi \mathbf{M}(0,1,4,7)$
4. $\quad F(A, B, C)=\Sigma \mathbf{m}(0,1,2,4,7)$
5. $\quad \mathbf{F}(\mathbf{A}, \mathbf{B}, \mathbf{C})=\Pi \mathbf{M}(0,1,2,4,5,7)$

Note: For both questions 13 and 14, use the truth table of the 3-to-8 decoder from Figure 10(a).
14. In Figure $10(b)$, function $\mathbf{G}(\mathbf{A}, \mathbf{B}, \mathbf{C})$ is equivalent to the function shown in:

1. Figure 11(a)
2. Figure 11(b)
3. Figure 11(c)
4. Figure 11(d)
5. Figure 11(e)


Figure 11 - Circuits for question 14.
15. How many 2-to-4 decoders are used to implement an 8-to-256 decoder?

Note: All the decoders have an enable input and all the inputs and outputs are un-inverted.
Hint: It is a hierarchical implementation with more than two levels of 2-to-4 decoders.

1. 5
2. 65
3. 69
4. 73
5. 85


Figure 12 - Circuit for questions 16 and 17. Note, $C_{0}$ is carry in and $C_{4}$ is carry out for the 4-bit adder.

Reminder: This adder and subtractor unit operates on 2's complement numbers and the $S$ input signal determines whether an addition or subtraction will occur.
16. In Figure 12, if $\mathbf{A}_{3} \mathbf{A}_{2} \mathbf{A}_{1} \mathbf{A}_{\mathbf{0}}=0011, \mathbf{B}_{3} \mathbf{B}_{2} \mathbf{B}_{1} \mathbf{B}_{\mathbf{0}}=0111$ and $\mathbf{S}=1$ then the output is:

1. $S u m=1100$ and $C_{4}=1$
2. $\operatorname{Sum}=1100$ and $C_{4}=0$
3. Sum=0100 and $C_{4}=1$
4. Sum=0100 and $C_{4}=0$
5. none of the above
6. In Figure 12, arithmetic overflow occurs for:
7. $\quad A_{3} A_{2} A_{1} A_{0}=0001, B_{3} B_{2} B_{1} B_{0}=1111$ and $\mathbf{S}=1$
8. $\quad \mathbf{A}_{3} \mathbf{A}_{2} \mathbf{A}_{1} \mathbf{A}_{\mathbf{0}}=0001, \mathbf{B}_{3} \mathbf{B}_{2} \mathbf{B}_{1} \mathbf{B}_{\mathbf{0}}=1111$ and $\mathbf{S}=0$
9. $\quad A_{3} A_{2} A_{1} A_{0}=1111, B_{3} B_{2} B_{1} B_{0}=0111$ and $\mathbf{S}=1$
10. $\quad \mathbf{A}_{3} \mathbf{A}_{2} \mathbf{A}_{1} \mathbf{A}_{\mathbf{0}}=1111, \mathbf{B}_{3} \mathbf{B}_{2} \mathbf{B}_{1} \mathbf{B}_{0}=0111$ and $\mathbf{S}=0$
11. none of the above
12. Consider a 1-bit adder cell with 3 inputs (1-bit operands $\mathbf{A}_{\mathbf{i}}, \mathbf{B}_{\mathbf{i}}$ and input carry $\mathbf{C}_{\mathbf{i}}$ ) and 2 outputs (1-bit sum $\mathbf{S}_{\mathbf{i}}$ and output carry $\mathbf{C}_{\mathbf{i + 1}}$ ). Then $\mathbf{C}_{\mathbf{i + 1}}$ can be implemented as shown in:
13. Figure 13(a)
14. Figure 13(b)
15. Figure 13(c)
16. Figure 13(d)
17. none of the above


Figure 13 - Implementations for output carry $C_{i+1}$ for question 18.
19. A 2-to-1 multiplexer illustrated in Figure 14(a) can be implemented using only 2-input NAND gates, as shown in Figure 14(b). Figure 14(c) gives the implementation of a function $f$ using an interconnection of three 2-to-1 multiplexers. It is assumed that the 2-input NAND gate's propagation delays from Low to High and High to Low are equal to 5 ns . Then the longest propagation delay from any input ( $\mathbf{d}_{3} \mathbf{d}_{2} \mathbf{d}_{1} \mathbf{d}_{0} \mathbf{S}_{2} \mathbf{S}_{1} \mathbf{S}_{\mathbf{0}}$ ) to the output (f) from Figure 14(c) is:

1. 25 ns
2. 30 ns
3. $\quad 35 \mathrm{~ns}$
4. 40 ns
5. 45 ns

Hint: Determine the propagation delay to the output of each NAND gate within every 2-to-1 multiplexer.

(a)

(b)

(c)

Figure 14 - Multiplexer circuits for question 19.
20. Consider the circuit from Figure 15(a). It is assumed that the 2-input NAND and 2-input XOR gates' propagation delays from Low to High and High to Low are equal to 5 ns. As shown in Figure 15(b), the input signal (input from Figure 15(a)) is a periodic signal with a frequency of 25 MHz . Then the output signal (output from Figure 15(a)) is a periodic signal and its frequency is:


Figure 15 - Circuit for question 20.

