Multimedia Communications



- In many lossy compression applications we want to represent source outputs using a small number of code words.
- Process of representing a large set of values with a much smaller set is called quantization
- Quantizer consists of two mappings: encoder mapping and decoder mapping
- Encoder divides the range of values that the source generates into a number of intervals
- Each interval is represented by a distinct codeword



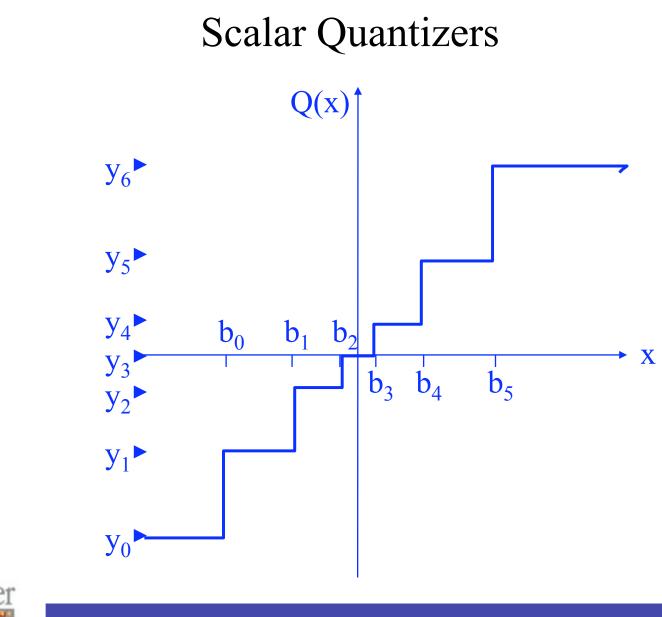
- All the source outputs that fall into a particular interval are represented by the codeword of that interval
- For every codeword generated by the encoder, the decoder generates a reconstruction value
- Because a codeword represents an entire interval, there is no way of knowing which value in the interval actually was generated by the source

Code	Output	
000	-3.5	
001	-2.5	
010	-1.5	
011	-0.5	
100	0.5	
101	1.5	
110	2.5	
111	3.5	



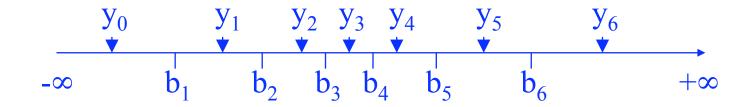
- Construction of intervals can be viewed as part of the design of the encoder
- Selection of reconstruction values is part of the design of the decoder
- The quality of the reconstruction depend on both intervals and reconstruction values
- The design is considered as a pair
- Design a quantizer: divide the input range into intervals, assign binary codes to these intervals, find reconstruction values
- Do all these while satisfying the rate-distortion criteria







- Source: random variable X with pdf of $f_x(x)$
- M: number of intervals
- b_i, i=0, 1,2, ..., M: M+1 end points of the intervals (decision boundary)
 - b₀ and b_M could be infinite
- y_i: M reconstruction level





• Distortion: mean squared quantization error

$$\delta_q^2 = E[(X - Q(X))^2] = \int_{-\infty}^{\infty} (x - Q(x))^2 f_X(x) dx$$

$$\delta_q^2 = \sum_{i=1}^M \int_{b_{i-1}}^{b_i} (x - y_i)^2 f_X(x) dx$$

- Rate: if fixed-length codewords are used to represent the quantizer output, the rate is given by: $R = \lceil \log_2 M \rceil$
- Selection of decision boundaries will not affect the rate



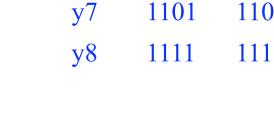
• If variable length codewords are used the rate will be:

$R = \sum_{i=1}^{M} l_i P(y_i)$	y1	1110	000
	y2	1100	001
b_i	v3	100	010

$$P(y_i) = \int_{b_{i-1}}^{b_i} f_x(x) dx$$
 y3 100
y4 00

$$R = \sum_{i=1}^{M} l_i \int_{b_{i-1}}^{b_i} f_x(x) dx$$
 y5
y6

Selection of decision boundaries will affect the rate



01

101

Code1 Code2

011

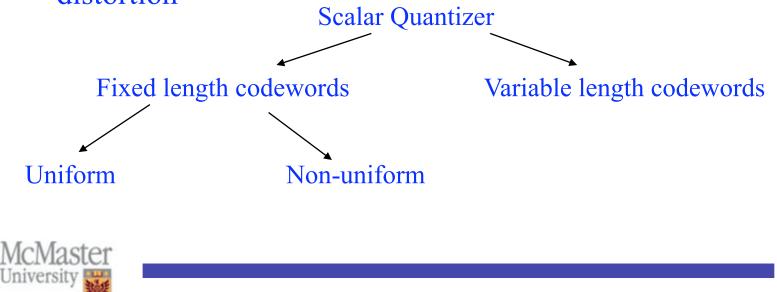
100

101

110



- Problem of finding optimum scalar quantizer:
- 1. given a distortion constraint $\sigma_q^2 \le D^*$ find the decision boundaries, reconstruction levels, and binary codes that minimize the rate or
- 2. given a rate constraint $R < R^*$ find the decision boundaries, reconstruction levels, and binary codes that minimize the distortion

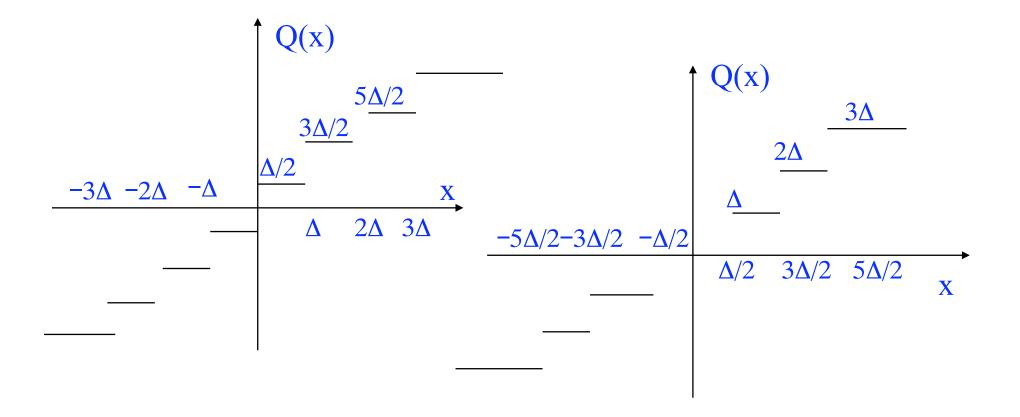


Uniform Quantizers

- All intervals are the same size except possibly for the two outer intervals (decision boundaries are spaced evenly)
- Reconstruction values are also spaced evenly with the same spacing as the decision boundaries
- In the inner intervals the reconstruction values are the midpoint of the intervals
- If zero is not a reconstruction level of the quantizer, it is called a midrise quantizer (M is even)
- If zero is a reconstruction level of the quantizer, it is called a midtread quantizer (M is odd)

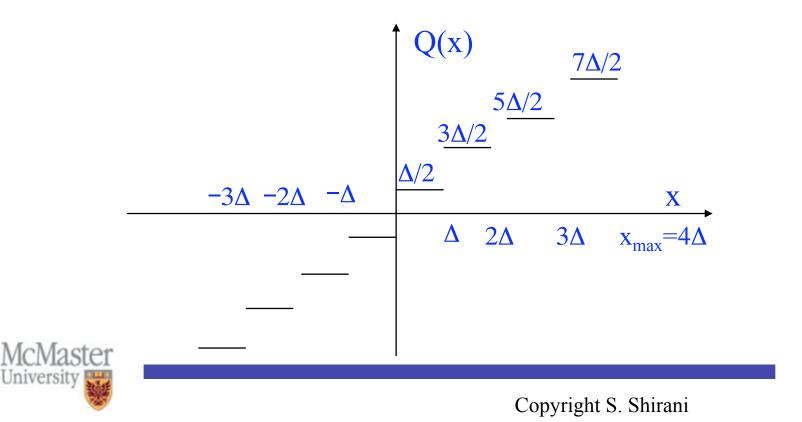


Uniform Quantizers



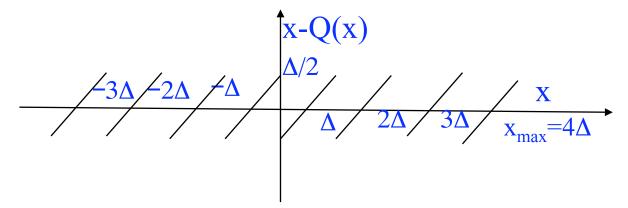


- Input: uniformly distributed in [-x_{max}, x_{max}]
- For an M-level uniform quantizer we divide $[-x_{max}, x_{max}]$ into M equally sized intervals each with a step size of $\Delta = \frac{2x_{max}}{M}$



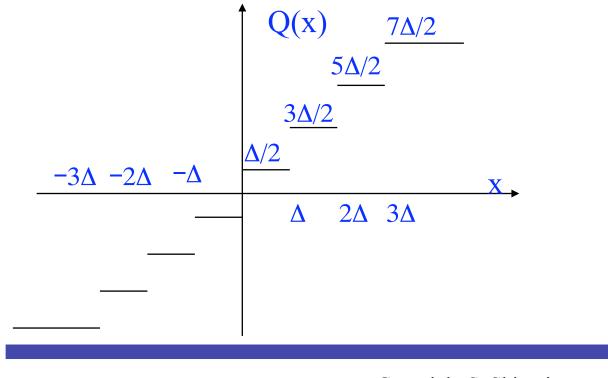
- For uniform distributions, quantization error q = Q(x) x of a uniform quantizer is uniformly distributed in $[-\Delta/2, +\Delta/2]$
- Then E[q] = 0, $\sigma_q^2 = E[q^2] = \Delta^2/12$.
- If the quantizer output is encoded using n bits per sample:

$$SNR(dB) = 10\log_{10}\frac{\sigma_x^2}{\sigma_q^2} = 10\log_{10}M^2 = 6.02\log_2 M = 6.02n$$





- When the distribution is not uniform, it is not a good idea to find the step size by dividing the range of inputs by the number of levels (sometimes the range is infinite)
- If pdf is not uniform, an optimal quantization step size can be found that minimizes D.





- Given the f_x(x) the value of step size can be calculated using numerical techniques
- Results are given for different values of M and different pdfs in table 9.3 of the book
- If input is unbounded, quantization error is no longer bounded
- In the inner intervals the error is bounded and is called granular error
- Unbounded error is called overload error

 $\begin{array}{c} x-Q(x) \\ \hline & \Delta/2 \\ \hline & X \\ \hline & granular noise \\ \end{array}$

- The pdf of most of non-uniform sources peaks at zero and decay away from the origin
- Overload probability is generally smaller than the probability of granular region
- Increasing step size will reduce overload error and increase granular error
- Loading factor is defined as the ratio of maximum value the input can take in the granular region to standard deviation
- Typical value for loading factor is 4

 $\begin{array}{c|c} x-Q(x) & & \text{overload noise} \\ \hline -3\Delta - 2\Delta - \Delta & & \Delta \\ \hline \Delta & 2\Delta - 3\Delta & \\ \hline x & \\ granular noise \end{array}$

- Mismatch effects: occur when the pdf of the input signal changes. The reconstruction quality degrades.
- Quantizer adaptation:
 - FORWARD: A block of data is processed, mean and variance sent as side info.
 - BACKWARD: based on quantized values, no side info. Jayant quantizer.
- Forward: a delay is necessary, size of block of data: if block is too large the adaptation process may not capture the changes in the input statistics, if too small transmission of side information adds a significant overhead



- If we study the input-output of a quantizer we can get an idea about the mismatch from the distribution of output values
- If Δ (quantizer step size) is smaller than what it should be, the input will fall in the outer levels of the quantizer an excessive number of times
- If Δ is larger than what it should be, the input will fall in the inner levels of the quantizer an excessive number of times
- Jayant quantizer: If the input falls in the outer levels, step size needs to be expanded, and if the input falls into inner levels, the step size needs to be reduced



- In Jayant quantizer expansion and contraction of the step size is accomplished by assigning a multiplier M_k to each interval.
- If the (n-1)th input falls in the kth interval, the step size to be used for the nth input is obtained by multiplying step size used for the (n-1)th input with M_k .
- Multiplier values for the inner levels in the quantizer are less than one and multiplier values for the outer levels are greater than one
- In math: $\Delta_n = M_{l(n-1)} \Delta_{n-1}$



- If input values are small for a period of time, the step size continues to shrink. In a finite precision system it would result in a value of zero.
- Solution: a minimum Δ_{\min} is defined and the step size cannot go below this value
- Similarly to avoid too large values for the step size we define a Δ_{max}



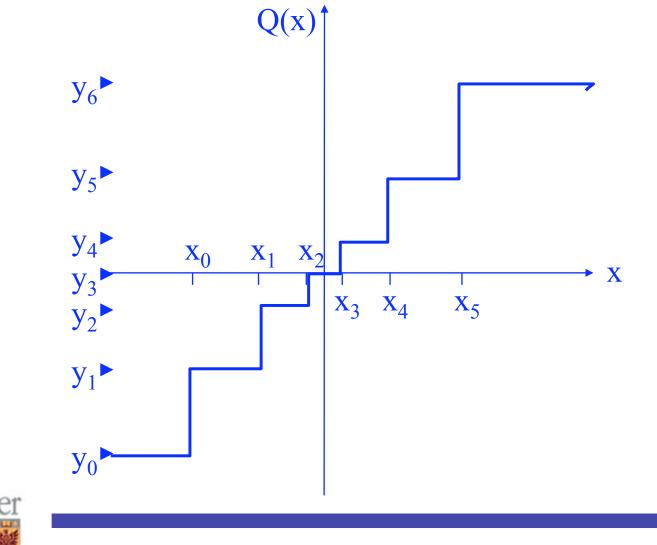
- How to choose the values of multipliers?
- Stability criteria: once the quantizer is matched to the input, the product of expansions and contractions are equal to one.

$$\begin{split} &\prod_{k=0}^{M} M_{k}^{n_{k}} = 1 \Longrightarrow \prod_{k=0}^{M} M_{k}^{n_{k}/N} = 1 \\ &\prod_{k=0}^{M} M_{k}^{P_{k}} = 1 \\ &M = \gamma^{l_{k}} \Longrightarrow \prod_{k=0}^{M} \gamma^{l_{k}P_{k}} = 1 \Longrightarrow \sum_{k=0}^{M} l_{k}P_{k} = 1 \end{split}$$



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Non-uniform Scalar Quantizers





Non-uniform Quantizers

- Non uniform quantizers attempt to decrease the average distortion by assigning more levels to more probable regions.
- For given M and input pdf, we need to choose $\{b_i\}$ and $\{y_i\}$ to minimize the distortion

$$\sigma^{2} = \sum_{i=1}^{M} \int_{b_{i-1}}^{b_{i}} (x - y_{i})^{2} f_{X}(x) dx$$
$$y_{i} = \frac{\int_{b_{i-1}}^{b_{i}} f_{X}(x) dx}{\int_{b_{i-1}}^{b_{i}} f_{X}(x) dx} \qquad b_{i-1} = \frac{y_{i-1} + y_{i}}{2}$$



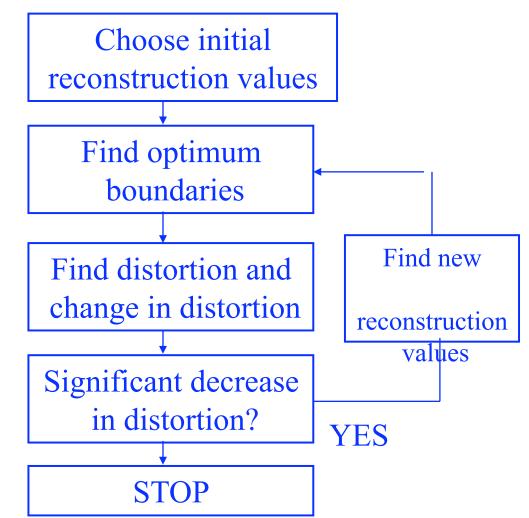
Non-uniform Quantizers

- The optimum value of reconstruction level depends on the boundary and the boundary depends on the reconstruction levels.
- Instead, it is easier to find these optimality conditions:
 - For a given partition (encoder), what is the optimum codebook (decoder)?
 - For a given codebook (decoder), what is the optimum partition (encoder)?

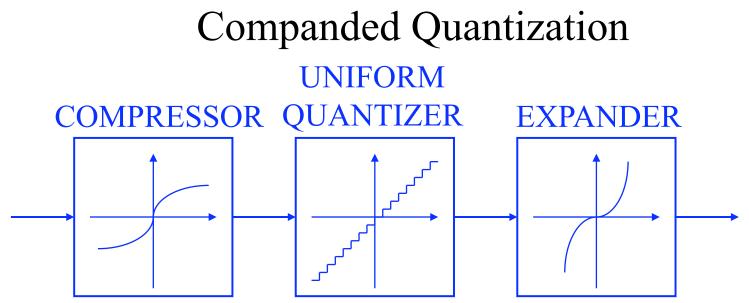


The Lloyd Algorithm

- It is difficult to solve both sets of equations analytically.
- An iterative algorithm known as the Lloyd algorithm solves the problem by iteratively optimizing the encoder and decoder until both conditions are met with sufficient accuracy.







- Instead of making the step size small for intervals in which the input lies with high probability, make these intervals large and use a uniform quantizer
- Equivalent to the a non-uniform quantizer.
- Example: the µ-law compander:

$$c(x) = x_{\max} \frac{\ln(1 + \mu \frac{|x|}{x_{\max}})}{\ln(1 + \mu)} \operatorname{sign}(x)$$



Entropy-Constrained Quantization

- Two approaches: 1) Keep the design of quantizer the same and entropy code the quantization output 2) Take into account the the selection of decision boundaries will affect the rate
- Joint optimization: entropy-constrained optimization. Minimized is prigged to the constraint

$$i P_{i} = \int_{b_{i-1}}^{b_{i}} f_{x}(x)dx \ln \frac{P_{k+1}}{P_{k}} = \lambda(y_{k+1} - y_{k})(y_{k+1} + y_{k} - 2b_{k}) y_{j} = \frac{\int_{b_{j-1}}^{b_{j}} xf_{x}(x)dx}{\int_{b_{j-1}}^{b_{j}} f_{x}(x)dx}$$



High-Rate Optimum Quantization

$$b_{k} = \frac{y_{k+1} + y_{k}}{2} - \frac{1}{2\lambda(y_{k+1} - y_{k})} \ln \frac{P_{k+1}}{P_{k}}$$
$$y_{j} = \frac{\int_{b_{j-1}}^{b_{j}} x f_{x}(x) dx}{\int_{b_{j-1}}^{b_{j}} f_{x}(x) dx}$$

- The above iterative method is called Generalized Lloyd algorithm.
- At high rates, the optimum entropy-constrained quantizer is the uniform quantizer!

• At high rates,
$$H(Q(X)) = h(X) - \log \frac{2x_{max}}{N}$$

